Efficient sequence assembly and variant calling using compressed data structures



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This dissertation is submitted for the degree of Doctor of Philosophy September 2012 I took a lengthy path to reach this point and my family supported me the entire way. I dedicate this work to them - thank you Mom, Dad, Calley and Kim.

Declaration

This dissertation describes work carried out from May 2009 to July 2012 under the supervision of Dr Richard Durbin at the Wellcome Trust Sanger Institute, while a member of Queens' College, Cambridge. This dissertation is the result of my own work and includes nothing which is the outcome of work done in collaboration except where specifically indicated in the text. The content in Chapter 2 was published in Simpson and Durbin [2010]. The content of Chapter 3 was published in Simpson and Durbin [2012].

This thesis does not exceed the length limit of 60,000 words as specified by the Biology Degree Committee.

Jared Thomas Simpson August 29, 2012

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Summary

De novo genome assembly is one of the most computationally demanding problems in genomics. In this thesis, I describe a collection of novel algorithms for performing *de novo* assembly using compressed data structures. First, I describe an algorithm to directly construct the assembly string graph - a model of overlap-based sequence assembly - using the compressed FM-index data structure. Previous algorithms for constructing the string graph required the intermediate step of building a full overlap graph, then removing transitive edges from the graph. My novel FM-index based algorithm does not require this time-consuming intermediate step. This algorithm allows fast and memory efficient overlap-based assembly. In Chapter 3, I extend my FM-index algorithms to build a space-efficient assembler for real sequencing data by designing error correction, read merging and scaffolding algorithms. Using these efficient algorithms I am able to reduce the memory requirement for assembling a human genome to 54GB.

In Chapter 4, I address the problem of detecting DNA sequence differences between two related genomes - the *variant calling* problem. Traditional approaches to variant calling align short sequence reads to a reference genome. While this approach is effective for simple differences, like isolated SNPs, it is more difficult to find complex changes like the insertion or deletion of sequence. My approach is based on analyzing the structure of an assembly graph built from the sequence data from multiple individuals. In Chapter 5, I apply this approach to real sequencing problems, including finding *de novo* mutations in the child of two parents, somatically acquired mutations in cancer and polymorphic variants present in a large human population.

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