

### **Outline of the Talk:**

- Global and Local Alignment
- □ Statistical significance of alignment
- Alignment method

### **Biological Motivation** Why We Need Sequence Alignment

### Inference of Homology

- Two genes are homologous if they share a common evolutionary history.
- Evolutionary history can tell us a lot about properties of a given gene
- Homology can be inferred from similarity between the genes
- Searching for Proteins with same or similar functions

# **Sequence Alignment**

#### **Global Alignment:**

#### Goal: How similar are two sequences $S_1$ and $S_2$

**Input:** two sequences  $S_1$ ,  $S_2$  over the same alphabet **Output:** two sequences  $S'_1$ ,  $S'_2$  of equal length ( $S'_1$ ,  $S'_2$  are  $S_1$ ,  $S_2$  with possibly additional gaps)

#### Example:

- $S_1 = \text{GCGCATGGATTGAGCGA}$
- $S_2$ = TGCGCCATTGATGACC
- A possible alignment:

 $S'_{1}$ = -GCGC-ATGGATTGAGCGA

S'\_= TGCGCCATTGAT-GACC--

# Sequence Alignment (cont)

**Local Alignment:** 

Goal: Find the pair of substrings in two input sequences which have the highest similarity

**Input:** two sequences  $S_1$ ,  $S_2$  over the same alphabet **Output:** two sequences  $S'_1$ ,  $S'_2$  of equal length ( $S'_1$ ,  $S'_2$  are substrings of  $S_1$ ,  $S_2$  with possibly additional gaps)

Example:

- $S_1$ = GCGCATGGATTGACCGA
- $S_2$ = TGCGCCATTGATGACC
- A possible alignment:

 $S'_1$  = ATTGA-G  $S'_2$  = ATTGATG

### **Global vs. Local Alignment**

- The <u>Global Alignment Problem</u> tries to find the longest path between vertices (0,0) and (n,m) in the edit graph.
- The Local Alignment Problem tries to find the longest path among paths between arbitrary vertices (*i*,*j*) and (*i*', *j*') in the edit graph.

### Global vs. Local Alignment (cont'd)

#### Global Alignment

# • Local Alignment—betten alignment to find conserved segment

tccCAGTTATGTCAGgggacacgagcatgcagagac

# **Local Alignment: Example**



### **Statistic Significance of Alignment**

We need to know how to evaluate the significance of the alignment. There are two scenarios:

First, the alignment indicates an evolutionary relationship between the sequences.

Second, the alignment is a chance occurrence. What answer is correct?

Here, the statistics are important to estimate of probability that the given alignment score might occur by chance.

#### E Value (E)

• E value (E) of an alignment score is the expected number of unrelated sequences in a database that would have a score at least as good.

# • Low E-values suggest that sequences are homologous.

If *E value*  $\leq$  0.02 sequence probably homologous

If *E value*  $\leq$  1 homology cannot be ruled out

If *E value* > 1 a match just by chance

- Statistical significance depends on both the size of the alignments and the size of the sequence database
  - Important consideration for comparing results across different searches
  - E-value increases as database gets bigger
  - E-value decreases as alignments get longer

**E value Measuring Alignment Significance** 

### P Value (P)

The *E-value* is not a probability; it's an expected value, i.e. the expected outcome.

Another criteria of the Alignment Significance is the probability that an alignment with this score could have arisen by chance - *p*-value:

E-value(S) =  $n \cdot p$ -value(S),

Here n is the number of sequences in the database, S.

The lower the p-value, the more likely it is that the alignment score is not by chance but was caused by alignment procedure.

For example, p = .01 means there is a 1 in 100 chance the result occurred by chance.

### **Methods of DNA Sequence Alignment**

- Dot matrix analysis
- The dynamic programming (DP) algorithm
  - Needleman-Wunsch Algorithm
  - Smith-Waterman Algorithm
- Suffix tree
- Hash table based algorithm
- Short read alignment tools

### **Dot Matrix Analysis**

• A dot matrix analysis is a method for comparing two sequences to look for possible alignment (Gibbs and McIntyre 1970)

• One sequence (A) is listed across the top of the matrix and the other (B) is listed down the left side

• Starting from the first character in B, one moves across the page keeping in the first row and placing a dot in many column where the character in A is the same

• The process is continued until all possible comparisons between A and B are made

- Any region of similarity is revealed by a diagonal row of dots
- Isolated dots not on diagonal represent random matches





## **The Needleman-Wunsch Algorithm**



### Smith-Waterman Algorithm

- Only works effectively when gap penalties are used
- In example shown
  - match = +1

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- mismatch = -1/3
- gap = -1+1/3k (k=extent
  of gap)
- Start with all cell values = 0
- Looks in subcolumn and subrow shown and in direct diagonal for a score that is the highest when you take alignment score or gap penalty into account

	С	Α	G	С	С	U	С	G	С	U	U	Α	G
Α	0.0	1.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	1.0	0.0
Α	0.0	1.0	0.7	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	1.0	0.7
U	0.0	0.0	0.8	0.3	0.0	0.0	0.0	0.0	0.0	1.0	1.0	0.0	0.7
G	0.0	0.0	1.0	0.3	0.0	0.0	0.7	1.0	0.0	0.0	0.7	0.7	1.0
С	1.0	0.0	0.0	2.0	1.3	0.3	1.0	0.3	2.0	0.7	0.3	0.3	0.3
С	1.0	0.7	0.0	1.0	3.0	1.7							
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 $H_{ij} = \max\{H_{i-1, j-1} + s(a_{ij}b_{j}), \max\{H_{i-k, j} - W_{k}\}, \max\{H_{i, j-1} - W_{l}\}, 0\}$ 

### **Bounded Dynamic Programming**

#### **Initialization:**

F(i,0), F(0,j) undefined for i, j > k

#### **Iteration:**

X<sub>M</sub>



$$\begin{split} F(i, j) &= max \, \begin{cases} F(i-1, j-1) + s(x_i, \, y_j) \\ F(i, j-1) - d, \, \text{if } j > i - k(N) \\ F(i-1, j) - d, \, \text{if } j < i + k(N) \end{cases} \end{split}$$

k(N)

Termination: same

Easy to extend to the affine gap case

# Suffix Tree Example





### Sequence Representation

Sequence S: 
$$(s_1s_2, ..., s_i, ..., s_m)$$
  $i = 1, 2, ..., m$   
K-tuple:  $(s_is_{i+1}...s_{i+k-1})$ 

Using two binary digits for each base, we may have the following representations:

For any of the m/k no-overlapping k-tuples in the sequence, an integer may be used to represent the k-tuple in a unique way

$$E = \sum_{i=1}^{2k} \beta_i 2^{i-1}$$
 with  $E_{\text{max}} = 2^{2k} - 1$ 

where  $\beta_i = 0$  or 1, depending on the value of the sequence base and  $E_{max}$  is the maximum **value** of the possible *E* values.

E	k-tuple	N <sub>i</sub>			Indi	ces and Of	fsets		
0	AA	1	2, 19						
1	AC	3	1, 9	2, 5	2, 11				
2	AG	2	1, 15	2, 35					
3	AT	2	2,13	3, 3					
4	CA	7	2, 3	2, 9	2, 21	2, 27	2, 33	3, 21	3, 2
5	СС	4	1, 21	2, 31	3, 5	3, 7			1
6	CG	1	1, 5						
7	СТ	6	1,23	2, 39	2,43	3, 13	3, 15	3, 17	
8	GA	4	1, 3	1, 17	2, 15	2, 25			
9	GC	0	11						
10	GG	5	1,25	1, 31	2, 17	2, 29	3, 1		
11	GT	6	1, 1	1, 27	1, 29	2, 1	2, 37	3, 19	
12	TA		3, 25						
13	TC	15	1, 7	1, 11	1, 19	2, 23	2,41	3, 11	
14	TG	3	1,13	2, 7	3, 9				

Hash Table





#### Mapping Score in ssaha2

Read mapping score is used to assess the repetitive feature of the read in the genome. With a maximum mapping score 50, we have:

$$S_{map} = \frac{10^* (30/R)(S_{max} - S_{max\,2})}{50} \quad if(S_{map} <= 50)$$
$$if(S_{map} > 50)$$

 $R = read length; S_{max}$  - maximum alignment score (smith-waterman) of the hits on genome;  $S_{max2}$  - second best alignment score of the hits on genome; Say you have one read of 30 bases which has a few hits on the genome: Best hit: exact match with  $S_{max}$  30; Second best hit: one base mismatch with  $S_{max2}$  29. The mapping score for this read is  $S_{map} = 10$ ;



### Genome Assembly using Solexa Short Reads Algorithms and Applications





# **Sequence Repeat Graph**









Genome/Chromosome

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l	2 1 IL2_33_8_1_571_251 max: 0 986799 1 12 0 0	NGACTAGTAAATTTGAAAAGTGGGCTGAGTGGCTGATGTTCTTGGATGCGGGGGN 12 A	
	$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	312         ARCHTCHGCCGCTTGGCCGACTTTTCGARTTASTRGT           ARCHTCHGCCACTCGGCCGACTTTTCGARTTASTRGT         43: 1225340595487         12824284357565         4962647         GANH           ARCHTCHGCCACTCGGCCGACTTTTCGARTTASTRGT         43: 1225340595487         12824284357565         4962647         GANH           ARCHARADTCHGLCHCTCTGGCCGACTTTGCARTTASTG         44: 1225340595487         12824284357565         4962647         GANH           ATCCHARANDTCHGLCHCTGGCCGACTTTTCGARTTASTG         44: 1225340595487         12824284357565         4962647         GANH           CHTCHARANDTCHGLCHCTGGCCGACTTTGCARTTASTG         44: 1225340595487         1282484357656         4962647         GANH           GCHTCHARANDTCHGLCHCTGGCCGACTTTGCART         45: 1225340595487         1282484357666         4962647         GANH           GCHTCHARANDTCHGCLGACTCHACTCHGCTCHAG         45: 1225340595487         128234835766         4962647         GANH           CCGCHTCHARANDTCHGGCLGACTCHCTTTCART         45: 122534059547         128248435766         4962647         GANH           CCGCHTCHARANDTCHGGCLGACTCHCTTTCART         45: 122534059547         128248435766         4962647         GANH           CCGCHTCHARANDTCHAGCLGACTCHGCUCTTTCA         45: 122534059547         128248435766         4962647         GANH           CCGCHTCHARANDTAGCLGATCHGCUCTTTCA         45: 122534059547	
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### Salmonella seftenberg Solexa Assembly from Pair-End Reads

Solexa reads:						
Number of reads:	6,00	6,000,000;				
Finished genome size:	~4.8	~4.8 Mbp;				
Read length:	2 <b>x</b> 37	2x37 bp;				
Estimated read coverage:	~92.	~92.5 X;				
Insert size:	170/	170/50-300 bp;				
Assembly features: - co	ntig stats					
	Solexa	454				
Total number of contigs:	75;	390				
Total bases of contigs:	4.80 Mbp	4.77 Mb				
N50 contig size:	139,353	25,702				
Largest contig:	395,600	62,040				
Averaged contig size:	63,969	12,224				
Contig coverage on genome:	<b>~99.8</b> %	<b>99.4</b> %				
Contig extension errors:	0					
Mis-assembly errors:	0	4				





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Genome Assembly - Normal Cell

Solexa reads:	
Number of reads:	557 Million;
Finished genome size:	3.0 GB;
Read length:	2x75bp;
Estimated read coverage:	~25x;
Insert size:	190/50-300 bp;
Number of reads clustered:	458 Million
Assembly features: - conti	a state

Assembly reactives concry	Stats
Total number of contigs:	1,020,346;
Total bases of contigs:	2.713 Gb
N50 contig size:	8,344;
Largest contig:	107,613
Averaged contig size:	2,659;
Contig coverage over the genome:	~90 %;
Mis-assembly errors:	?

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